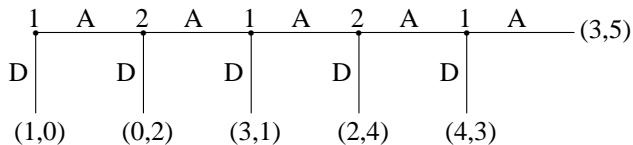


# Centipede Game



- Play this as a fun game...

# Computing Subgame Perfect Equilibria

Idea: Identify the equilibria in the bottom-most trees, and adopt these as one moves up the tree

```

function BACKWARDINDUCTION (node  $h$ ) returns  $u(h)$ 
if  $h \in Z$  then
   $\sqsubset$  return  $u(h)$ 
 $best\_util \leftarrow -\infty$ 
forall  $a \in \chi(h)$  do
   $util\_at\_child \leftarrow \text{BACKWARDINDUCTION}(\sigma(h, a))$ 
  if  $util\_at\_child_{\rho(h)} > best\_util_{\rho(h)}$  then
     $\sqsubset$   $best\_util \leftarrow util\_at\_child$ 
return  $best\_util$ 

```

- $util\_at\_child$  is a vector denoting the utility for each player
- the procedure doesn't return an equilibrium strategy, but rather labels each node with a vector of real numbers.
  - This labeling can be seen as an extension of the game's utility function to the non-terminal nodes
  - The equilibrium strategies: take the best action at each node.

# Computing Subgame Perfect Equilibria

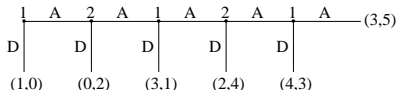
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```

- For zero-sum games, BACKWARDINDUCTION has another name: the **minimax** algorithm.
  - Here it's enough to store one number per node.
  - It's possible to speed things up by **pruning** nodes that will never be reached in play: "alpha-beta pruning".

# Backward Induction



- What happens when we use this procedure on Centipede?
  - In the only equilibrium, player 1 goes down in the first move.
  - However, this outcome is Pareto-dominated by all but one other outcome.
- Two considerations:
  - practical: human subjects don't go down right away
  - theoretical: what should you do as player 2 if player 1 doesn't go down?
    - SPE analysis says to go down. However, that same analysis says that P1 would already have gone down. How do you update your beliefs upon observation of a measure zero event?
    - but if player 1 knows that you'll do something else, it is rational for him not to go down anymore... a paradox
    - there's a whole literature on this question