

Pithy Quote

If there is intelligent life on other planets, in a majority of them, they would have discovered correlated equilibrium before Nash equilibrium.

– Roger Myerson

Examples

- Consider again Battle of the Sexes.
 - Intuitively, the best outcome seems a 50-50 split between (F, F) and (B, B) .
 - But there's no way to achieve this, so either someone loses out (unfair) or both players often miscoordinate
- Another classic example: traffic game

| | | |
|-----------|--------------|-------------|
| | <i>go</i> | <i>wait</i> |
| <i>go</i> | $-100, -100$ | $10, 0$ |
| <i>B</i> | $0, 10$ | $-10, -10$ |

Intuition

- What is the natural solution here?

Intuition

- What is the natural solution here?
 - A traffic light: a fair randomizing device that tells one of the agents to go and the other to wait.
- Benefits:
 - the negative payoff outcomes are completely avoided
 - fairness is achieved
 - the sum of social welfare exceeds that of any Nash equilibrium
- We could use the same idea to achieve the fair outcome in battle of the sexes.
- Our example presumed that everyone perfectly observes the random event; not required.
- More generally, some random variable with a commonly known distribution, and a private signal to each player about the outcome.
 - signal doesn't determine the outcome or others' signals; however, correlated

Formal definition

Definition (Correlated equilibrium)

Given an n -agent game $G = (N, A, u)$, a **correlated equilibrium** is a tuple (v, π, σ) , where v is a tuple of random variables $v = (v_1, \dots, v_n)$ with respective domains $D = (D_1, \dots, D_n)$, π is a joint distribution over v , $\sigma = (\sigma_1, \dots, \sigma_n)$ is a vector of mappings $\sigma_i : D_i \mapsto A_i$, and for each agent i and every mapping $\sigma'_i : D_i \mapsto A_i$ it is the case that

$$\begin{aligned} \sum_{d \in D} \pi(d) u_i(\sigma_1(d_1), \dots, \sigma_i(d_i), \dots, \sigma_n(d_n)) \\ \geq \sum_{d \in D} \pi(d) u_i(\sigma_1(d_1), \dots, \sigma'_i(d_i), \dots, \sigma_n(d_n)) . \end{aligned}$$

Existence

Theorem

For every Nash equilibrium σ^ there exists a **corresponding correlated equilibrium** σ .*

- This is easy to show:
 - let $D_i = A_i$
 - let $\pi(d) = \prod_{i \in N} \sigma_i^*(d_i)$
 - σ_i maps each d_i to the corresponding a_i .
- Thus, correlated equilibria always exist

Remarks

- Not every correlated equilibrium is equivalent to a Nash equilibrium
 - thus, correlated equilibrium is a **weaker notion** than Nash
- Any **convex combination of the payoffs** achievable under correlated equilibria is itself realizable under a correlated equilibrium
 - start with the Nash equilibria (each of which is a CE)
 - introduce a second randomizing device that selects which CE the agents will play
 - regardless of the probabilities, no agent has incentive to deviate
 - the probabilities can be adjusted to achieve any convex combination of the equilibrium payoffs
 - the randomizing devices can be combined